

CSC236 *Intro. to the Theory of Computation*

Lecture 9: Finite State Automata

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Course page:

<http://www.cdf.toronto.edu/~csc236h/fall/index.html>

Section page:

http://www.cdf.toronto.edu/~csc236h/fall/amir_lectures.html

review

❖ so far

- different flavour of proofs and their application in cs

❖ in particular, recently

we saw tools useful toward

- **proof**: if a program is **semantically** correct

let's reword it:

- **recognize**: if a program is **semantically** correct

❖ next: finite state machines/automata

- tools useful to **recognize** if a program is **syntactically** correct
- **and ...**

Examples 83, 84

❖ identifiers

- e.g., a *letter* followed by a *digit*

- more practical ones

Example 85

❖ Python-like float

Example 86

- ❖ strings with an odd number of a 's (and any number of b 's)

Finite State Automaton definition

- ❖ is a 5-tuple $M = (Q, \Sigma, \delta, q_0, F)$
 - Q is the set of states, which is finite & non-empty
 - Σ is the alphabet, which is finite & non-empty
 - $\delta : Q \times \Sigma \rightarrow Q$ is the transition function
 - $q_0 \in Q$ is the start state
 - $F \subseteq Q$ is the set of accept states

- ❖ Then, $L(M)$ is a language
 - that machine M accepts,
 - i.e., set of all strings that machine M accepts

Example 86 revisited

- ❖ devise a machine that only accepts strings with an odd number of a 's. $\Sigma = \{a, b\}$

Examples 85, 84, 83 revisited

- ❖ **85.** devise a *machine* that accepts strings representing a float number a . $\Sigma = \{0..9, +, -, .\}$
- ❖ **84.** devise a *machine* that accepts identifiers
 $\Sigma = \{0..9, a..z, _ \}$
- ❖ **83.** devise a *machine* that accepts simple identifiers (length 2, first character a letter). $\Sigma = \{0..9, a..z\}$

notes

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