CSC 165

bounds and problems
week 9, lecture 2
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A2: - use CDF submission facility early to
work out glitches

- Use any application where you can
type input + produce PDF

Proposed tutorials:

something from A2

Assignment 2 wants you to consider the negation of: $\exists x \in \mathbb{Q}^{\geq 0}, \exists \epsilon \in \mathbb{Q}^{+}, \forall \delta \in \mathbb{Q}^{+}, \exists y \in \mathbb{Q}^{\geq 0}(|x-y| < \delta \land |2x-2y| \geq \epsilon)$ $\forall \chi \in \mathbb{Q}^{\geq 0}, \forall \xi \in \mathbb{Q}, \exists \delta \in \mathbb{Q}^{+}, \forall \chi \in \mathbb{Q}^{\geq 0}, \forall \chi \in \mathbb{Q}^{\geq 0$

There are a couple of ways to negate the conjunction. The one we have a straight-forward proof technique for is implication.

bounded below

recall

Notice that the definition of big-Omega differs in just one character from big-Oh:

$$oxed{\Omega(g)} = \{f: \mathbb{N} \mapsto \mathbb{R}^{\geq 0} \mid \exists c \in \mathbb{R}^+, \exists B \in \mathbb{N}, orall n \in \mathbb{N}, n \geq B \Rightarrow f(n) \geq cg(n)\}$$

The rôle of B is, as with big-Oh, to act as a breakpoint, so comparisons don't have to start at the origin

The rôle of c is to scale g down below f.

If you're proving $f \in \Omega(g)$, you get to choose c and B to suit your proof.

one last bound

It often happens that functions are bounded above and below by the same function. In other words, $f \in \mathcal{O}(G) \land f \in \Omega(g)$. We combine these two concepts into $f \in \Theta(g)$.

$$\Theta(g) = \{f: \mathbb{N} \mapsto \mathbb{R}^{\geq 0} \mid \exists c_1 \in \mathbb{R}^+, \exists c_2 \in \mathbb{R}^+, \exists B \in \mathbb{N}, orall n \in \mathbb{N}, n \geq B \Rightarrow c_1 g(n) \leq f(n) \leq c_2 g(n) \}$$

You might want to draw pictures, and conjecture about values of c_1, c_2, B for $f = 5n^2 + 15$ and $g = n^2$.

 $\mathcal{A} = \mathcal{E}_{f}: \mathbb{N} \to \mathbb{R}^{3}$ How about: $g \in \Omega(f) \Leftrightarrow f \in \mathcal{O}(g)$ $\forall f, ge f, ge SC(f) \Rightarrow fe O(g)$ Cossume $f,g \in f \neq generic$ Assume $g \in SZ(f) \neq antecedent$ Then $F \subset ER^+, B \in IN, \forall n \in N, n \geq B \Rightarrow f(n) \geq cg(n)$ Choose $C_1 \in \mathbb{R}^1$, $B_1 \in IN$, $\forall n \in IN$, $n \ge B_1 \Rightarrow g(n) \ge C_1 \in I$.

Let $C_2 = \frac{1/c_1}{B_1}$. Then $C_2 \in \mathbb{R}^+$ Let $B_2 = B_1$ Bz= B1 o Then Bz & IN. assume n e IN. Bosse Then $g(n) \ge C_1 f(n)$ # from antecedent Then $g(n) \ge C_1 f(n)$ # $B_2 = B_1$ Thus Czg(n) = CzG, f(n) = f(n), # Cz= = Then 12B, slide 5

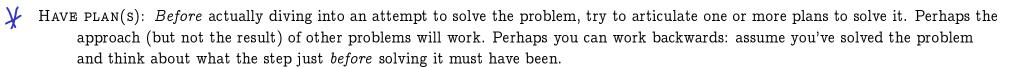
scratch

problem solving

Often problem-solving techniques aren't taught in mathematically-oriented courses. Rather than expect you to have picked them up on your, we offer some approaches that sometimes work.

We start from the approach elaborated by George Polya in How to Solve it

Understand the problem: What's given (input), what's required (output)? Can you represent either input or output in ways that seem clearer — possibly symbols or diagrams? Make sure you can state the problem clearly in your own words.



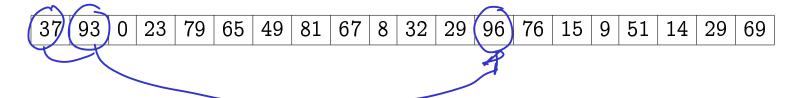
TRY OUT A PLAN: If it works, try to verify it. If it doesn't work, think about where you got stuck (even write down what blocked you).

Did you have a plan B in the previous step?

LOOK BACK: If you get some results, look back on what worked and what didn't. Can you extend or generalize the problem into a new problem?

today's problem

You can get a handout to work on in groups. Here's the gist of the problem: Suppose you're given the list:



Can you find one or more longest non-decreasing sequences? For example, 37, 93, 96 is a sequence that's non-decreasing, but you can easily find longer ones.

Sequences are ordered, but need not be contiguous