CSC 165
proof by cases
week 7, lecture 3
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proof by cases

mathematic tes pre requisites

You can prove by induction (CSC236) that:

$$orall m\in \mathbb{N}, orall n\in \mathbb{N}, n
eq 0 \Rightarrow \exists !q\in \mathbb{N}, \exists !r\in \mathbb{N}, m=qn+r ext{ and } n>r\geq 0$$

 \exists ! is a compact way of saying there exists exactly one.

q and r are called the quotient and remainder, respectively. We also denote

$$r ext{ by } m ext{ mod } n = r$$

integer division - some as 90 for non-neg m,n "implementation dependent"

A consequence is that any natural number n has a remainder of either 0, 1, or 2 after division by 3. What possible remainders are there for perfect squares

$$0^{2} \mod 3 = 0$$
 $1^{2} \mod 3 = 1$
 $2^{2} \mod 3 = 1$
 $3^{2} \mod 3 = 0$

$$4^{2}$$
 mod $3 = 1$
 5^{2} mod $3 = 1$
 6^{2} mod $3 = 6$

Prove: $\forall n \in \mathbb{N}, n^2 \bmod 3 \neq 2$

I think you'll need cases for different possible results of $n \mod 3$

assume n e IN # generic
Then n mod 3 E \ \ 20,1,23 # defn of mod 3 Case!, assume n mod 3 = 0 Then $\exists q' \in IN$, n = 3q' + 0 # defin of n mod <math>3 = 0Pick $q \in \mathbb{N}$, n = 3q # Since it exist. Then $n^2 = 3(3q^2)$ # algebra Then $3q \in \mathbb{N}$, $n^2 = 3q' + 0$ # $3q^2$, Since $3, q \in \mathbb{N}$ and $2 \in \mathbb{N}$, $n^2 = 3q' + 0$ # $3q^2$, Since $3, q \in \mathbb{N}$ and Then 1 mod 3 = 0 = 2 Case 2, assume n mod 3=1 Then I g'EN, n=3g'+1 # Jefn of n mod 3=1

Then 12 mon 3 = 1 + 2

scratch

Case 3, assume n mod 3 = 2 Then $\exists q' \in IN, n = 3q' + 2$ Puck $q \in IN, n = 3q + 2$ Then n' = qq' + 12q + 4 $= 3(3q^2 + 4q + 1) + 1$ # Since It exists. # algebra. # more algebra # 9 = (39,+49+1) # e IN, suro IN # closed X, + # and 3,4,1,9 EIN So $\exists q \in \mathbb{N}, n^2 = 3q+1$ Then n 2 mod 3 = 1 + 2

Conclude FNEIN, no mod 3 # 2# Since true
in all possible coses.

slide 4

more modular arithmetic

Performing arithmetic on integers, and then taking the remainder, can be interchanged with taking the remainder and then performing the arithmetic. Number theory and cryptography use many such techniques.

Example: $\forall m \in \mathbb{N}, \forall n \in \mathbb{N}, (m \bmod 7 = 3 \land n \bmod 7 = 4) \Rightarrow mn \bmod 7 = 5$

An important step is to remember that $m \mod 7 = 3$ means there are natural numbers q and p such that m = 7q + 3.

Ymin EN

Prove: $\forall m \in \mathbb{N}, \forall n \in \mathbb{N}, (m \bmod 7 = 3 \land n \bmod 7 = 4) \Rightarrow mn \bmod 7 = 5$

assume m, n e IN # generic pair in IN. Assume (m mod 7 = 3 1 n mod 7 = 4) # antecedent

scratch

Then MN mod 7 = 5Conclude $4 \text{ m, n } \in \text{IN}$, $(\text{m mod } 7 = 3 \text{ } \Lambda \text{ n mod } 7 = 4)$ $\implies \text{mn mod } 7 = 5$

structural transformations

Pages 48–49 of the notes list some standard transformations used in proofs. These have the virtue that many proof practitioners will recognize them when you use them.

Work back over proofs to find:

- implication introduction
- universal introduction
- existential introduction
- existential elimination