### CSC148 winter 2018

binary trees week 8

```
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```

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### Outline

general trees continued...

binary trees

traversals

binary search trees

## queues, stacks, recursion

You may have noticed in the last slide there were no recursive calls, and a queue was used to process a recursive structure in level order.

Careful use of a stack allows you to process a tree in preorder.



### tree inheritance issues

one approach to BinaryTree would be to make it a subclass of Tree, but there are some design considerations

- ▶ any client code that uses Tree would be required not to violate the branching factor (2) of BinaryTree
- one way to achieve this would be to make Tree immutable: make sure there is no way to change children or value, and then have subclasses that might be mutable

we will take a different approach: a completely separate BinaryTree class





#### **BTNode**

Change our generic Tree design so that we have two named children, left and right, and can represent an empty tree with None

# special methods...

We'll want the standard special methods:

#### contains

True

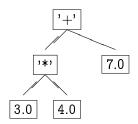
```
you've implemented contains on linked lists, nested Python lists, general
Trees before; implement this function, then modify it to become a method

def contains(node: BTNode, value: object) -> bool:
    """
    Return whether tree rooted at node contains value.

>>> contains(None, 5)
False
>>> contains(BTNode(5, BTNode(7), BTNode(9)), 7)
```

# arithmetic expression trees

Binary arithmetic expressions can be represented as binary trees:





# evaluating a binary expression tree

- ▶ there are no empty expressions
- ▶ if it's a leaf, just return the value
- otherwise...
  - ▶ evaluate the left tree
  - evaluate the right tree
  - combine left and right with the binary operator

Python built-in eval might be handy.





#### inorder

#### A recursive definition:

- visit the left subtree inorder
- visit this node itself
- visit the right subtree inorder

The code is almost identical to the definition.



## preorder

- visit this node itself
- ▶ visit the left subtree in preorder
- visit the right subtree in preorder



## postorder

- ▶ visit the left subtree in postorder
- visit the rightsubtree in postorder
- visit this node itself



### level order

- ▶ visit root
- visit root's children
- ▶ visit root's grandchildren
- ▶ visit root's greatgrandchildren
- **.**..

### definition

Add ordering conditions to a binary tree:

- data are comparable
- ▶ data in left subtree are less than node.data
- data in right subtree are more than node.data



# why binary search trees?

Searchs that are directed along a single path are efficient:

- ▶ a BST with 1 one has height 1
- ▶ a BST with 3 nodes may have height 2
- ▶ a BST with 7 nodes may have height 3
- a BST with 15 nodes may have height 4
- ▶ a BST with n nodes may have height  $\lceil \lg n \rceil$ .



### bst\_contains

If node is the root of a "balanced" BST, then we can check whether an element is present in about  $\lg n$  node accesses.

```
def bst_contains(node: BTNode, value: object) -> bool:
    """
    Return whether tree rooted at node contains value.

Assume node is the root of a Binary Search Tree

>>> bst_contains(None, 5)
    False
    >>> bst_contains(BTNode(7, BTNode(5), BTNode(9)), 5)
    True
    """
# use BST property to avoid unnecessary searching
```





### mutation: insert

```
.....
Insert data in BST rooted at node if necessary, and return new root
Assume node is the root of a Binary Search Tree.
>>> b = BTNode(8)
>>> b = insert(b, 4)
>>> b = insert(b, 2)
>>> b = insert(b, 6)
>>> b = insert(b, 12)
>>> b = insert(b, 14)
>>> b = insert(b, 10)
>>> print(b)
        14
    12
        10
8
        6
```

def insert(node: BTNode, data: object) -> BTNode: