

CSC148 winter 2018

recursive structures

week 7

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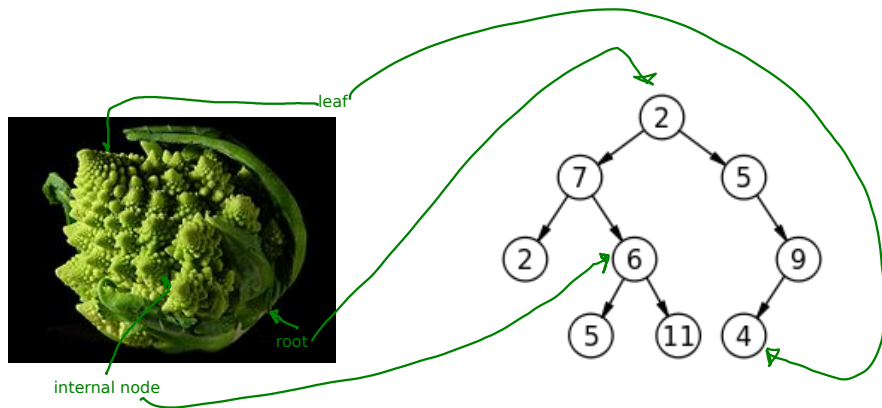
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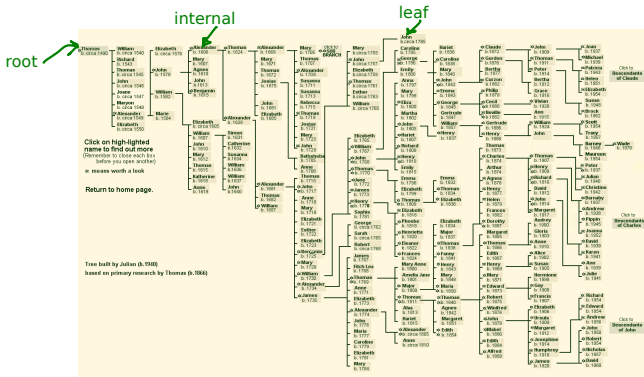


recursion, natural and otherwise



structure to organize information

patriarchal family tree...



terminology

- ▶ set of **nodes** (possibly with values or labels), with directed **edges** between some pairs of nodes
- ▶ One node is distinguished as **root** edges are directed away from root
- ▶ Each non-root node has exactly one parent.
...root has zero!
- ▶ A **path** is a sequence of nodes n_1, n_2, \dots, n_k , where there is an edge from n_i to n_{i+1} . The **length** of a path is the number of edges in it
paths lead away from the root because edges do...
there is a path of length 0 from a node to itself...
- ▶ There is a unique path from the root to each node. In the case of the root itself this is just n_1 , if the root is node n_1 .
unique: one and only one, aka exactly one...
- ▶ There are no **cycles** no paths that form loops.
compare with csc165 definition of tree...



more terminology

- ▶ **leaf:** node with no children
- ▶ **internal node:** node with one or more children

- ▶ **subtree:** tree formed by any tree node together with its descendants and the edges leading to them.



- ▶ **height:** $1 +$ the maximum path length from the root to some leaf.

NB: Some texts use a different definition. Be sure you know which!

- ▶ **depth:** length of a path from root to a node is the node's depth.

*root - depth = 0
children of root - depth = 1*

- ▶ **arity, branching factor:** maximum number of children for any node.



general tree implementation

```
class Tree:
    """
    A bare-bones Tree ADT that identifies the root with the entire tree
    """
    def __init__(self, value: object=None,
                  children: List[Tree]=None) -> None:
        """
        Create Tree self with content value and 0 or more children
        """
        self.value = value
        # copy children if not None
        self.children = children[:] if children are not None else []
```

This approach avoids the dreaded mutable default parameter...

shallow copy...



general form of recursion:

if $\langle \text{condition to detect a base case} \rangle$:

Figure out 1 or more base cases and how to deal with them...

$\langle \text{do something without recursion} \rangle$

else: # $\langle \text{general case} \rangle$

Anything that isn't a base case involves 1 or more recursive calls. Assume they work correctly and figure out how to combine them...

$\langle \text{do something that involves recursive call(s)} \rangle$



how many leaves?

```
def leaf_count(t: Tree) -> int:
```

```
    """
```

```
    Return the number of leaves in Tree t.
```

```
>>> t = Tree(7)
```

```
>>> leaf_count(t)
```

```
1
```

```
>>> t = descendants_from_list(Tree(7),  
                                descendants_from_list used to build example (below)
```

```
                                [0, 1, 3, 5, 7, 9, 11, 13], 3)
```

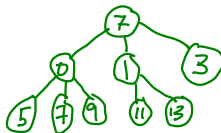
```
>>> leaf_count(t)
```

```
6
```

```
    """
```

```
    Base case when t is a leaf... how many leaves is that?
```

```
    General case: sum up the leaves in t's children...
```



height of this tree?

```
def height(t: Tree):  
    """  
    Return 1 + length of longest path of t.  
  
    >>> t = Tree(13)  
    >>> height(t)  
    1  
    >>> t = descendants_from_list(Tree(13),  
                                   [0, 1, 3, 5, 7, 9, 11, 13], 3)  
    >>> height(t)  
    3  
    """
```

1 more edge than the maximum height of a child, except
what do we do if there are no children?

base case when t is a leaf... how tall is a leaf?

general case: maximum height of children + 1...



arity, or branching factor

```
def arity(t: Tree) -> int:
```

```
    """
```

```
    Return the maximum branching factor (arity) of Tree t.
```

```
>>> t = Tree(23)
```

23

```
>>> arity(t)
```

```
0
```

```
>>> tn2 = Tree(2, [Tree(4), Tree(4.5), Tree(5), Tree(5.75)])
```

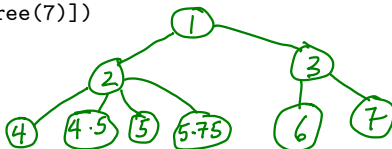
```
>>> tn3 = Tree(3, [Tree(6), Tree(7)])
```

```
>>> tn1 = Tree(1, [tn2, tn3])
```

```
>>> arity(tn1)
```

```
4
```

```
    """
```



filesystem example

```
from os import scandir, path
from tree import Tree
```

Try this on your computer...

```
def path_to_tree(path_name: str) -> Tree:
    """
    Return a Tree representing filesystem starting from pathname.
    """
    return Tree((path_name, [f.name for f in scandir(path_name)]),
                [path_to_tree(path.join(path_name, f.name))
                 for f in scandir(path_name)
                 if f.is_dir()]])
```



pass in a function

```
def count_if(t: Tree, p: Callable[[object], bool]) -> int:
    """
    Return number of values in Tree t that satisfy predicate p(value).
```

Assume predicate p is defined on t's values

```
>>> def p(v): return v > 4
>>> t = descendants_from_list(Tree(0),
                                [1, 2, 3, 4, 5, 6, 7, 8], 3)
```

```
>>> count_if(t, p)
```

```
4
```

```
>>> def p(v): return v % 2 == 0
```

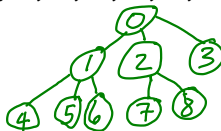
```
>>> count_if(t, p)
```

```
5
```

```
"""
```

base case: t is a leaf... return 0 if its value doesn't satisfy the predicate, otherwise 1...

general case: t has some children... return all the counts for t's children
plus 1 if t's value satisfies the predicate...



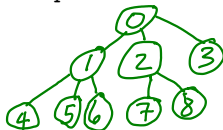
list the leaves

```
def list_leaves(t: Tree) -> int:
    """
    Return list of values in leaves of t.

    >>> t = Tree(0)
    >>> list_leaves(t)
    [0]
    >>> t = descendants_from_list(Tree(0),
                                   [1, 2, 3, 4, 5, 6, 7, 8], 3)
    >>> list_ = list_leaves(t)
    >>> list_.sort() # so list_ is predictable to compare
    >>> list_
    [3, 4, 5, 6, 7, 8]
    """
```

base case, t is a leaf: list of t's value

general case: concatenate lists of t's childrens leaves...



pre-order visit

```
def preorder_visit(t: Tree, act: Callable[[Tree], Any]) -> None:
    """
```

Visit each node of Tree `t` in preorder, and act on the nodes as they are visited. [see page 19 for a diagram...](#)

```
>>> t = descendants_from_list(Tree(0),
                                [1, 2, 3, 4, 5, 6, 7], 3)
```

```
>>> def act(node): print(node.value)
```

```
>>> preorder_visit(t, act)
```

```
0
```

```
1
```

```
4
```

```
5
```

```
6
```

```
2
```

```
7
```

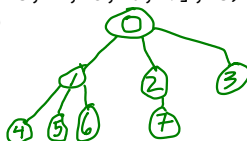
```
3
```

```
"""
```

```
act(t)
```

```
for c in t.children:
```

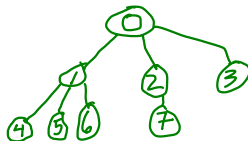
```
    preorder_visit(c, act)
```



postorder

```
def postorder_visit(t: Tree, act: Callable[[Tree], Any] -> None:
    """
    Visit each node of t in postorder, and act on it when it is visited

    >>> t = descendants_from_list(Tree(0),
                                     [1, 2, 3, 4, 5, 6, 7], 3)
    >>> def act(node): print(node.value)
    >>> postorder_visit(t, act)
    4
    5
    6
    1
    7
    2
    3
    0
    """
```



See code on page 15, and change the order slightly...



levelorder

```
def levelorder_visit(t: Tree, act: Callable[[Tree], Any]) -> None:  
    """
```

Visit every node in Tree *t* in level order and act on the node as you visit it.

```
>>> t = descendants_from_list(Tree(0),  
                                [1, 2, 3, 4, 5, 6, 7], 3)
```

```
>>> def act(node): print(node.value)
```

```
>>> levelorder_visit(t, act)
```

0

1

2

3

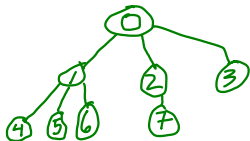
4

5

6

7

```
    """
```



Use either a queue, or a helper method that visits all the nodes at a given depth...



queues, stacks, recursion

You may have noticed in the last slide there were no recursive calls, and a **queue** was used to process a recursive structure in level order.

Careful use of a **stack** allows you to process a tree in preorder

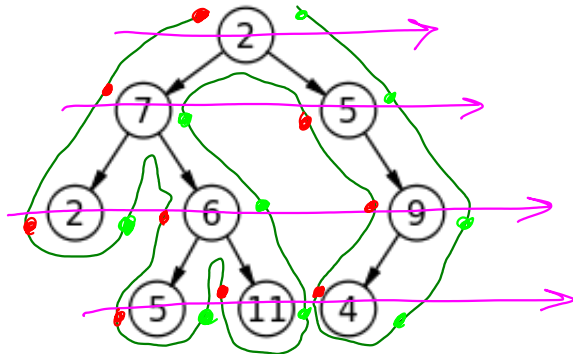


preorder tracing...

preorder

postorder

levelorder



notes...