Hashing*

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^{*}Some slides are taken from Danny Heap and Bogdan Simion, also youtube channel Kevin Drumm

Agenda

- Talk about hashing
- Course Evaluation

Lists and linear search

Data Structure	search()	insert()	delete()
List	O(n)	O(n)	O(n)

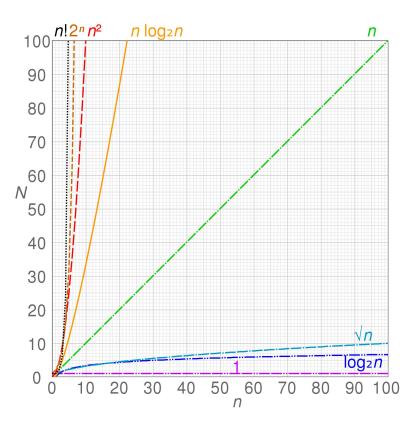
Ordered Lists and Binary search

Data Structure	search()	insert()	delete()
List	O(n)	O(n)	O(n)
Ordered List	O(lg n)	O(lg n)	O(lg n)

Demo

Onto Pycharm

Complexity Comparisons



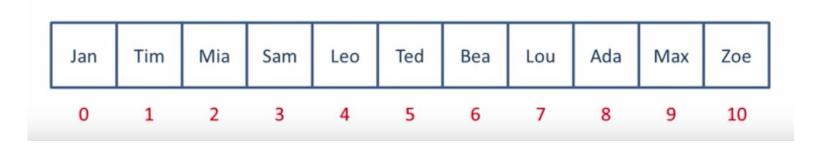
How dictionary offers Constant time search?

By a technique called hashing

Example: hashing

What's the time complexity of:

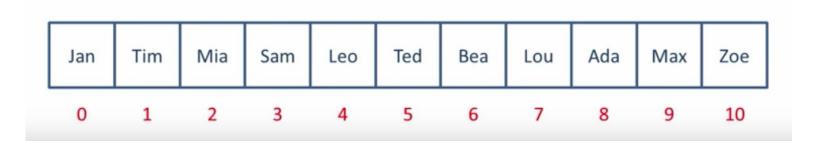
- 1) Finding the string 'Ada' in this list?
- 2) What if you knew the index of 'Ada'?



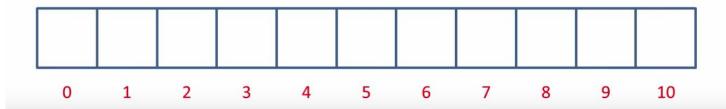
Example: hashing

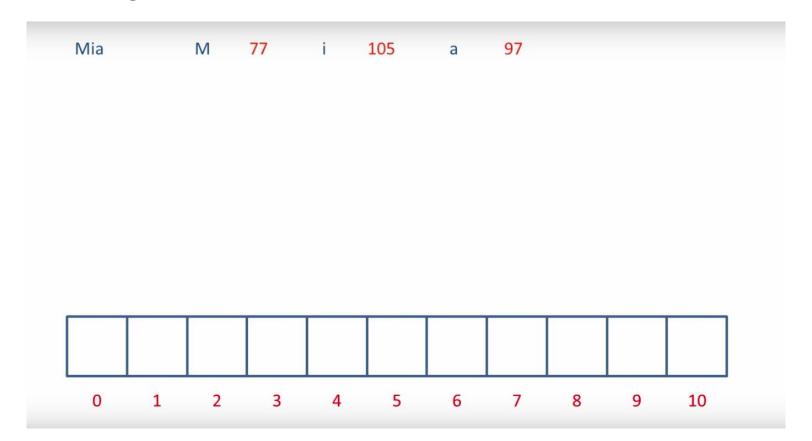
What's the time complexity of:

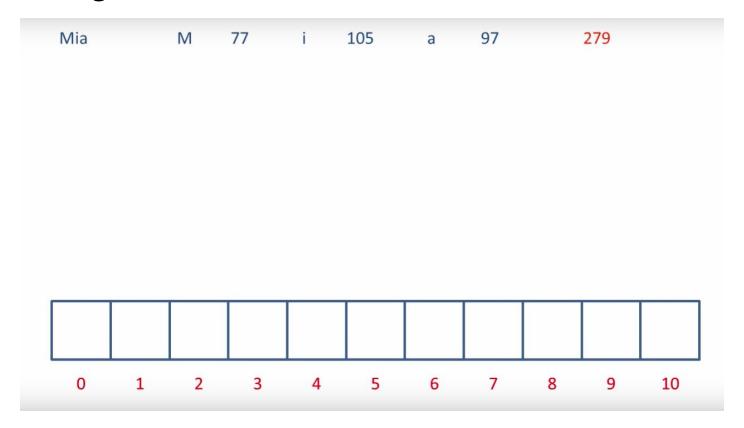
- 1) Finding the string 'Ada' in this list?
- 2) What if you knew the index of 'Ada'?
 - a) But how do you know the index of 'Ada' in the list?

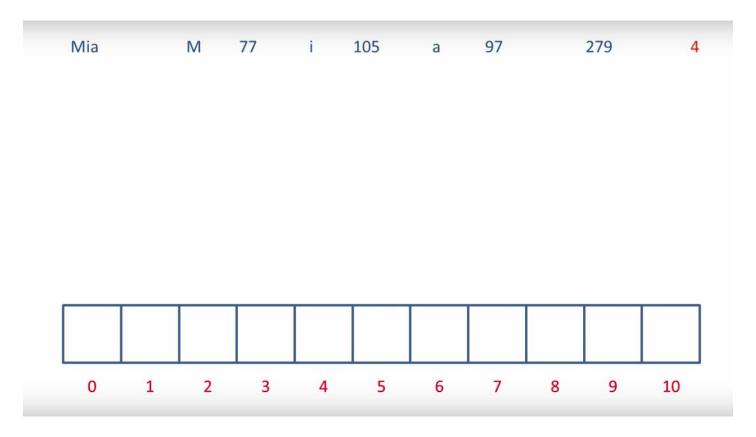


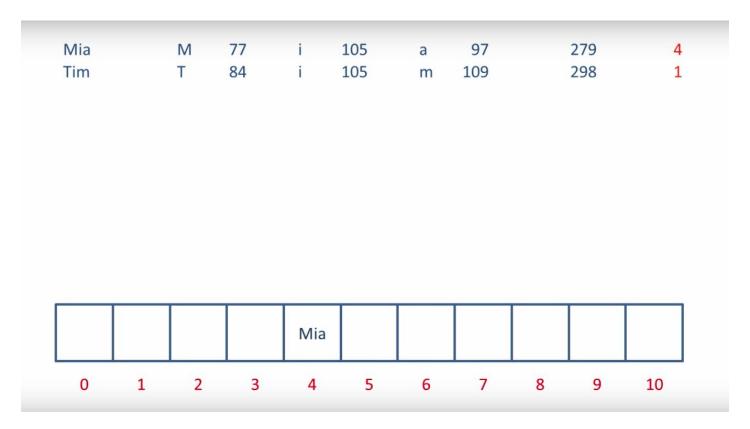
Mia M i a

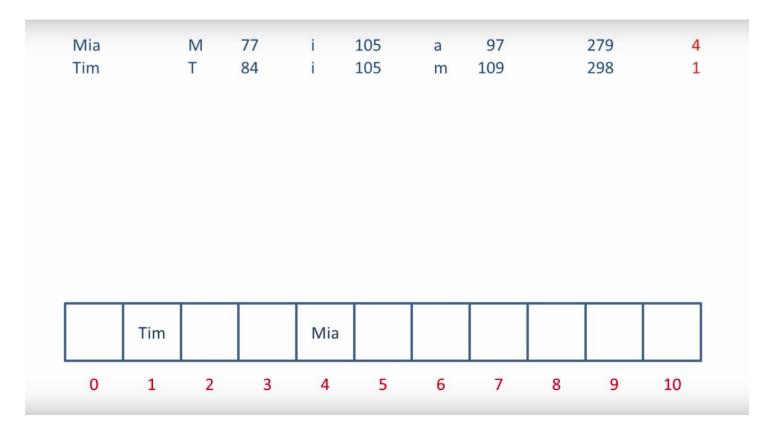


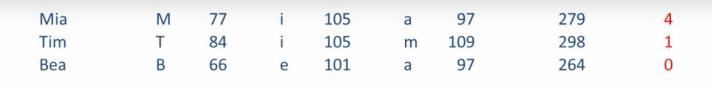


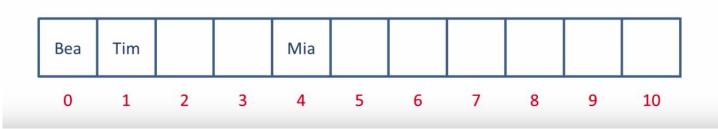






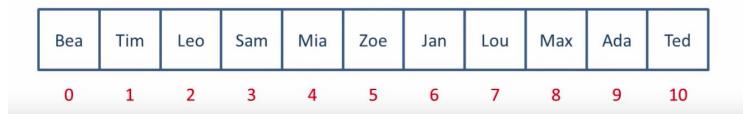




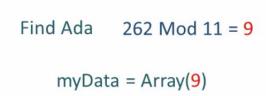


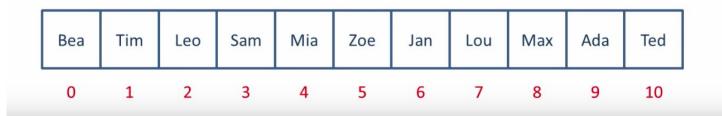
0	1	2	3	4	5	6	7	8	9	10
Bea	Tim	Leo	Sam	Mia	Zoe	Jan	Lou	Max	Ada	Ted
Ted		Т	84	е	101	d	100		285	10
Max		M	77	а	97	X	120		294	8
Lou		L	76	0	111	u	117		304	7
Sam		S	83	а	97	m	109		289	3
Leo		L	76	е	101	0	111		288	2
Ada		Α	65	d	100	a	97		262	9
Jan		J	74	а	97	n	110		281	6
Zoe		Z	90	0	111	е	101		302	5
Bea		В	66	е	101	a	97		264	0
Tim		Т	84	i	105	m	109		298	1
Mia		M	77	i	105	a	97		279	4

Index number = sum ASCII codes Mod size of array



Search





Key, value pairs



Key Value pairs



Key:Value Pairs

Keys

Bea 27/01/1941 English Astronomer	Tim U8/U6/1955 English Inventor	31/12/1945 American Mathematician	Sam 27/04/1791 American Inventor	Mia 20/02/1986 Russian Space Station	Zoe 19/06/1978 American Actress	Jan 13/02/1956 Polish Logician	Lou 27/12/1928 French Biologist	Max 23/04/1898 German Physicist	Ada 18/12/1815 English Mathematician	American
0	1	2	3	4	5	6	7	8	9	10

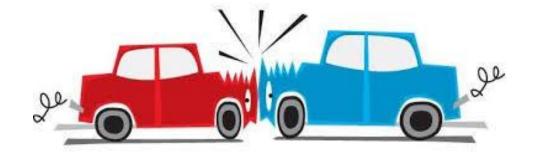
Values

Hashing Algorithm

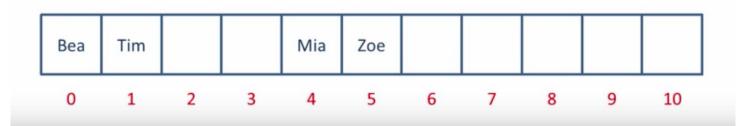
- Calculation applied to a key to transform it into an address
- For numeric keys, divide the key by the number of available addresses, n, and take the remainder

address = key Mod n

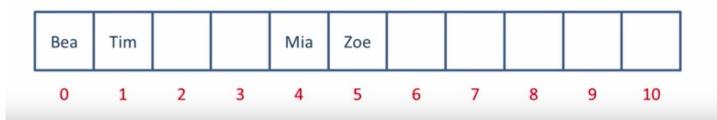
- For alphanumeric keys, divide the sum of ASCII codes in a key by the number of available addresses, n, and take the remainder
- Folding method divides key into equal parts then adds the parts together
 - The telephone number 01452 8345654, becomes 01 + 45 + 28 + 34 + 56 + 54 = 218
 - Depending on size of table, may then divide by some constant and take remainder



Mia	M	77	i	105	а	97	279	4
Tim	T	84	i	105	m	109	298	1
Bea	В	66	е	101	а	97	264	0
Zoe	Z	90	0	111	е	101	302	5

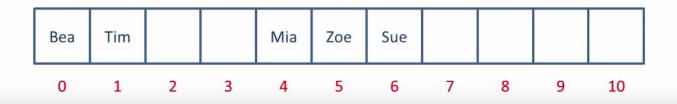


Mia	M	77	i	105	а	97	279	4
Tim	Т	84	i	105	m	109	298	1
Bea	В	66	е	101	a	97	264	0
Zoe	Z	90	0	111	е	101	302	5
Sue	S	83	u	117	е	101	301	4



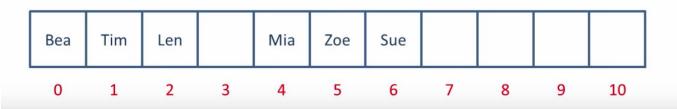
Collisions: Linear Probing

Mia	M	77	i	105	а	97	279	4
Tim	Т	84	i	105	m	109	298	1
Bea	В	66	е	101	а	97	264	0
Zoe	Z	90	0	111	е	101	302	5
Sue	S	83	u	117	е	101	301	4

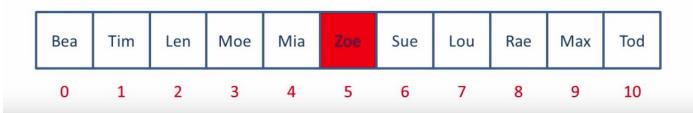


Collisions: Linear Probing

Mia	M	77	i	105	а	97	279	4	
Tim	Т	84	i	105	m	109	298	1	
Bea	В	66	е	101	а	97	264	0	
Zoe	Z	90	0	111	е	101	302	5	
Sue	S	83	u	117	e	101	301	4	
Len	L	76	е	101	n	110	287	1	

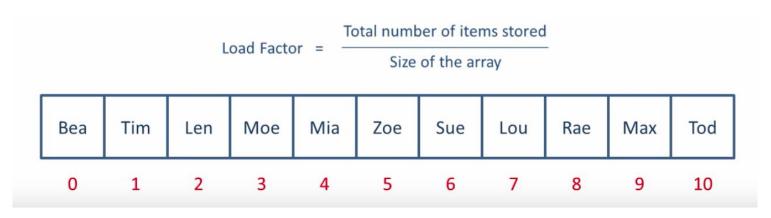


Collisions: Linear Probing

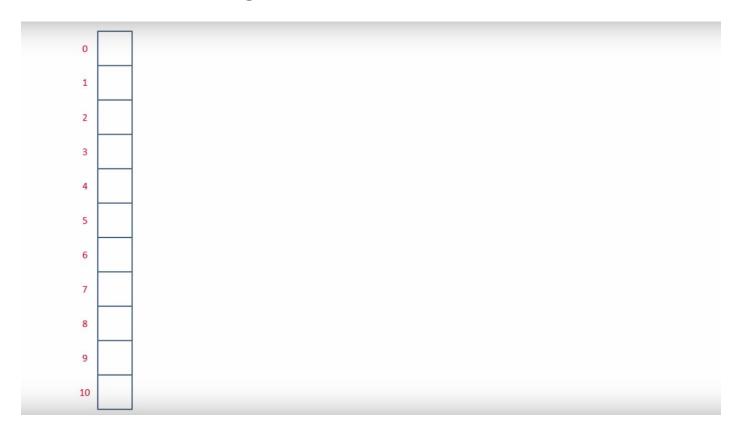


More data → more probability for collision

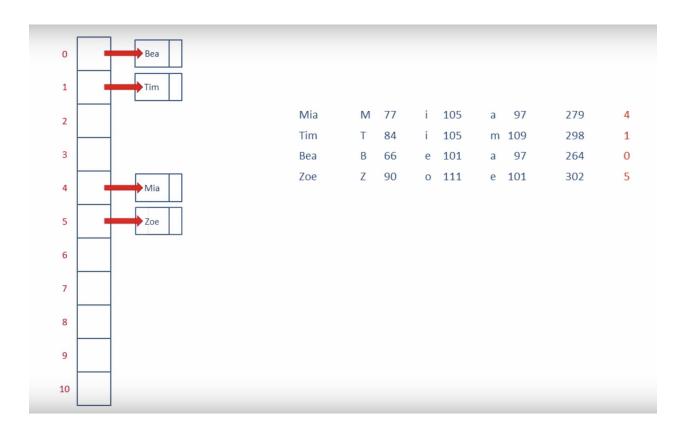
Make the hash table big enough as compared to the number of keys

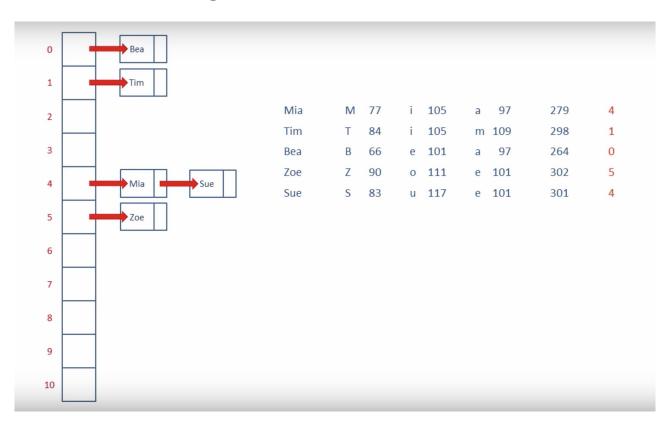


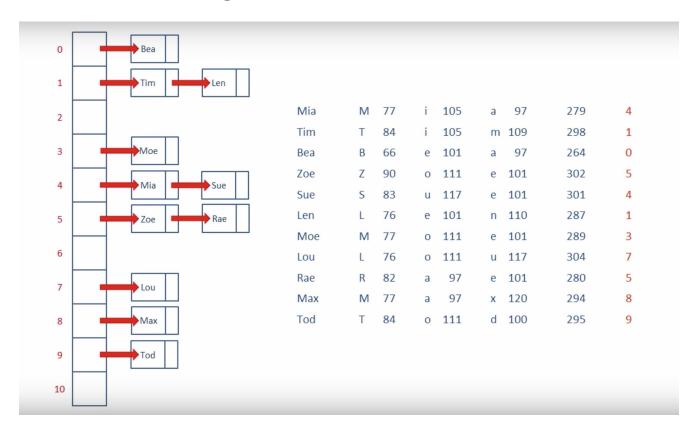
Resize the array based on certain threshold of *Load Factor*

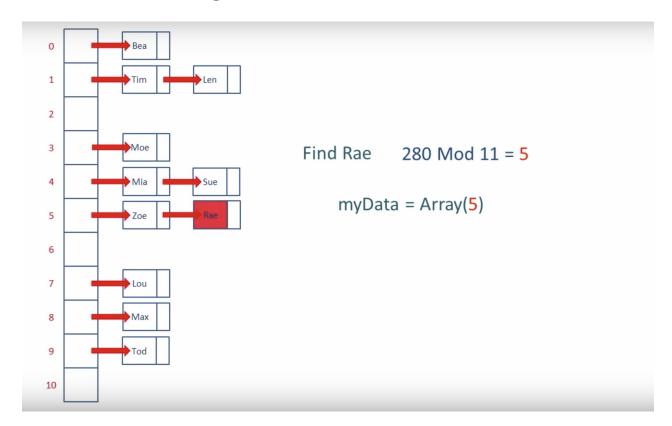












Collison Resolution

- Open addressing
 - Linear probing
 - Plus 3 rehash
 - Quadratic probing (failed attempts)²
 - Double hashing
- Closed addressing

How probable are collisions?

Highly Probable!

Birthday paradox:

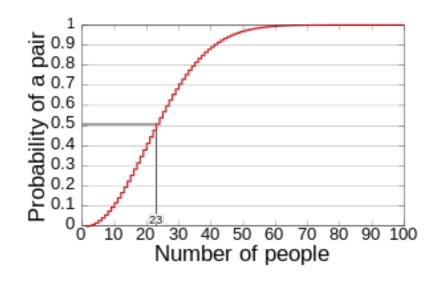
Given a room of 'n' people, at least how many people you require to find among them any two people with matching birthdays with 50% probability?

the mathematics is a bit counter-intuitive... the probability of a non-collision for 23 birthdays is:

$$p = \frac{366}{366} \times \frac{365}{366} \times \cdots \times \frac{344}{366} \approx 0.493$$

How probable are collisions?

Birthday Paradox



For 32-bit hash values: only 77,000 elements are required for significant risk (50%) of collision

Objectives of Hash Function

- Minimize collisions
- Uniform distribution of hash values
- Easy to calculate

Next Lecture

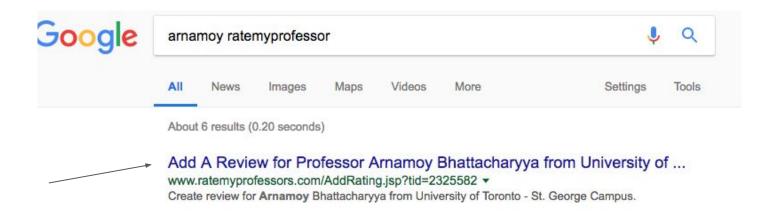
Implementing dictionary using Chaining

Course Evaluations

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Developing our own Python Dictionary

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