CSC148 Intro. to Computer Science

Lecture 9: BST (insert, delete)

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Course webpage:

http://www.cdf.toronto.edu/~csc148h/winter

Last week

- Binary trees (branch factor =2)
- Depth-first traversal
 - inorder, preorder, and postorder
- Breadth-first traversal
 - level-order
- Binary Search Trees
- Today
 - More on BST
 - insert
 - · delete

Binary Search Trees

- Add ordering conditions to a binary tree:
 - data are comparable
 - data in left subtree are less than node.data
 - data in right subtree are more than node.data

Binary Search Trees

- a BST with I node has height I
- a BST with 3 nodes may have height 2
- a BST with 7 nodes may have height 3
- a BST with 15 nodes may have height 4
- * a BST with n nodes may have height $\lceil \lg n \rceil$

- * if the BST is "balanced", then we can check whether an element is present in about $\lg n$ node accesses
 - This is significantly faster than a linear search: O(n)

bst_contains

```
def bst_contains(node, value):
    Return whether tree rooted at node contains value.
    Assume node is the root of a Binary Search Tree
    @param BinaryTree|None node: node of a Binary Search Tree
    Oparam object value: value to search for
    @rtype: boo7
    >>> bst_contains(None, 5)
    Fa1se
    >>> bst_contains(BinaryTree(7, BinaryTree(5), BinaryTree(9)), 5)
    True
    if node is None:
        return False
    elif value < node.data:
        return bst_contains(node.left, value)
    elif value > node.data:
        return bst_contains(node.right, value)
    else:
        return True
```

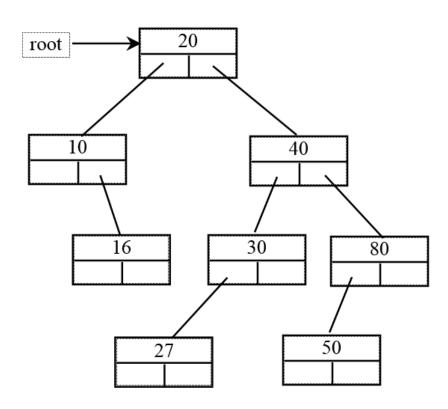
bst insert

```
def insert(node, data):
    Insert data in BST rooted at node if necessary, and return new root.
    Assume node is the root of a Binary Search Tree.
    @param BinaryTree node: root of a binary search tree.
    Oparam object data: data to insert into BST, if necessary.
    >>> b = BinaryTree(5)
    >>> b1 = insert(b, 3)
    >>> print(b1)
    <BLANKLINE>
    return_node = node
    if not node:
        return_node = BinaryTree(data)
    elif data < node.data:</pre>
        node.left = insert(node.left, data)
    elif data > node.data:
        node.right = insert(node.right, data)
    else: # nothing to do
        pass
    return return node
```

bst_insert

Let's trace it for a few examples:

bst_delete

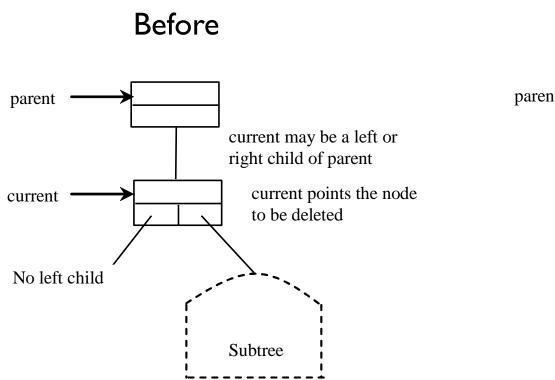


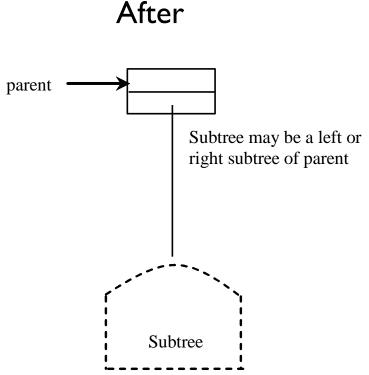
bst_delete

- First locate the node that contains the element and also its parent node.
- Let <u>current</u> point to the node that contains the element in the tree and <u>parent</u> point to the parent of the <u>current</u> node.
- There are two cases to consider ...

Case I: The current node has no left child

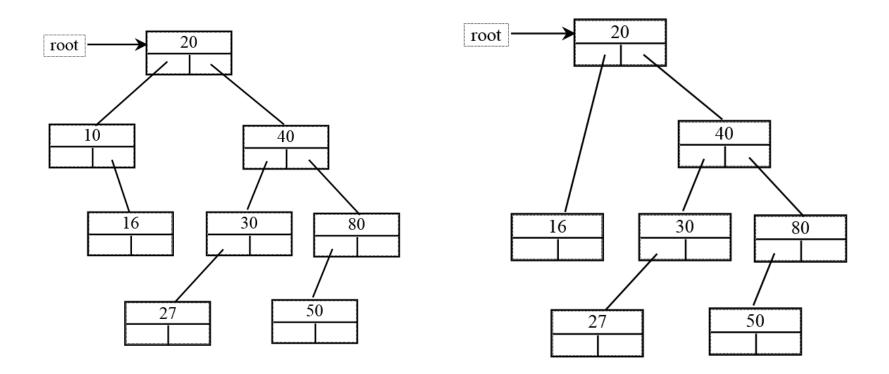
Simply connect the <u>parent</u> with the right child of the current node.





Example for Case 1. Deleting node 10

Connect the parent of node 10 with the right child of node 10.



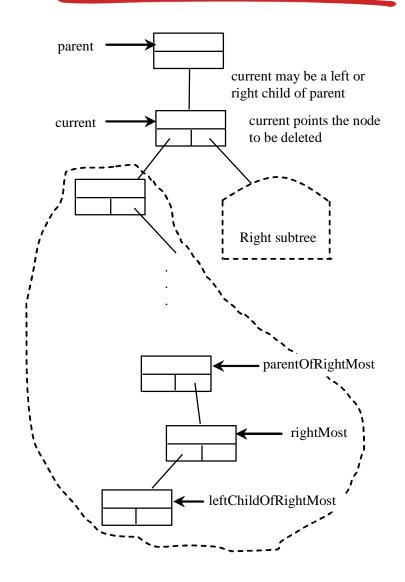
Case 2: The current node has a left child.

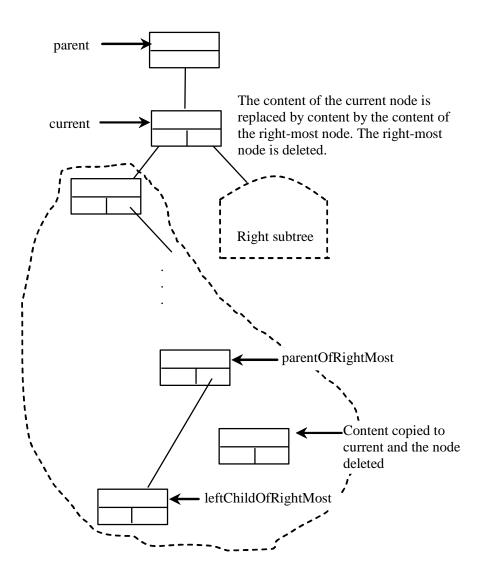
- Let <u>right_most</u> point to the node that contains the largest element in the left subtree of the <u>current</u> node.
- Let <u>parent_of_right_most</u> point to the parent node of the <u>right_most</u> node.

Then:

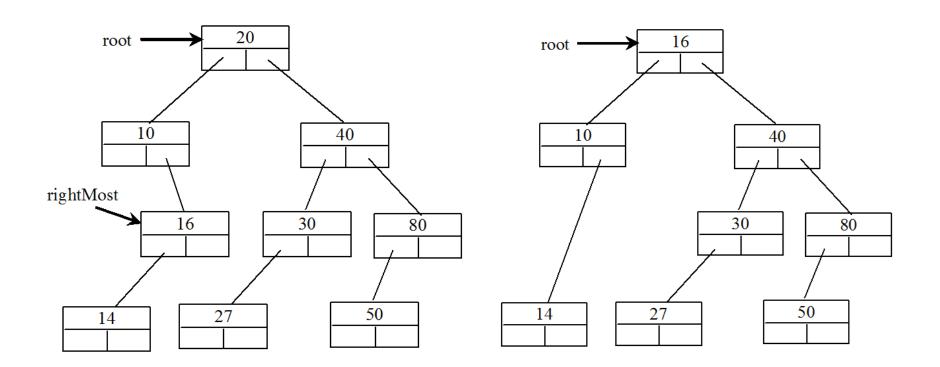
- I. Replace the element value in the <u>current</u> node with the one in the <u>right_most</u> node,
- 2. Connect the <u>parent_of_right_most</u> node with the left child of the <u>right_most</u> node.

Case 2 (diagram)

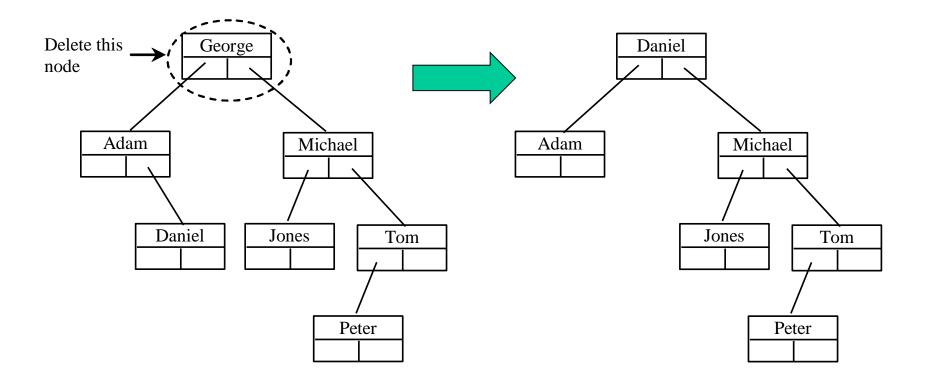




Example for Case 2. Deleting node 20

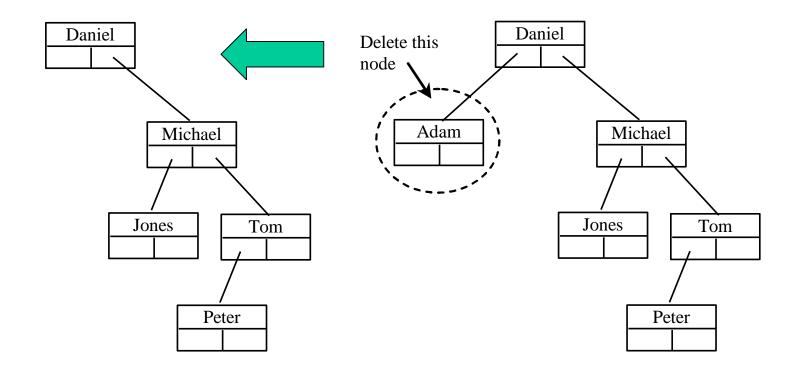


More Examples



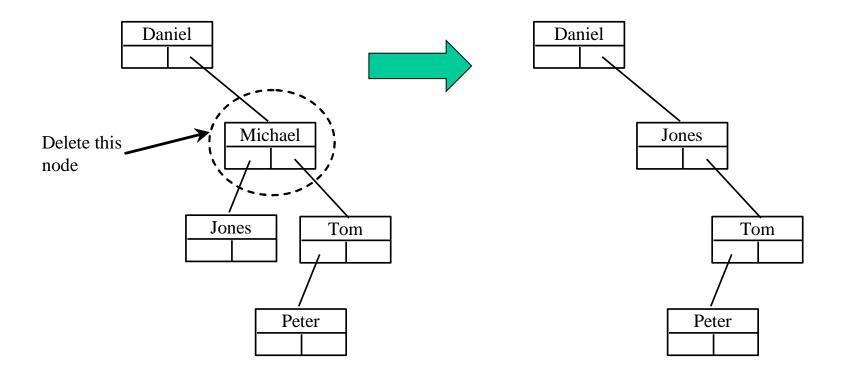
Case I or 2? 2

More Examples



Case I or 2?

More Examples



Case I or 2? 2

bst delete

 First locate the nodes that contain the element and its parent. Call them <u>current</u> and <u>parent</u>.

```
parent = None
current = root
while current is not None and current.data != data:
    if data < current.data:</pre>
        parent = current
        current = current.left
    elif data > current.data:
        parent = current
        current = current.right
    else: pass # Element is in the tree pointed at by current
if current is None: return False # Element is not in the tree
```

Case I: bst delete

```
# Case 1: current has no left child
if current left is None:
    # Connect the parent with the right child of the
      current node
    # Special case, assume the node being deleted is at
      root
    if parent is None:
        current = current.right
    else:
        # Identify if parent left or parent right should
          he connected
        if data < parent.data:</pre>
            parent.left = current.right
        else:
            parent.right = current.right
else:
    # Case 2: The current node has a left child
```

Case II: bst delete

```
# Locate the rightmost node in the left subtree of
# the current node and also its parent
parent_of_right_most = current
right_most = current.left
while right_most.right is not None:
    parent_of_right_most = right_most
    right_most = right_most.right # Keep going to the right
# Replace the element in current by the element in rightMost
current.element = right_most.element
# Eliminate rightmost node
 if parent_of_right_most.right == right_most:
    parent_of_right_most.right = right_most.left
else:
    # Special case: parent_of_right_most == current
    parent_of_right_most.left = right_most.left
return True # Element deleted successfully
```

Summary

- Homework:
 - In Slides 10 and 12,
 - replace every left with right, every right with left, and also largest with smallest.
 - And, implement the method.
- Next Week:
 - How bst_delete can be written recursively?