

# CSC148 winter 2015

linked structures

week 7

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# Outline

## Assignment 2

## binary trees

## traversals

## binary *search* trees



## tippy and minimax

This continues the game-playing framework of Assignment 1, adding a new game and a new strategy:

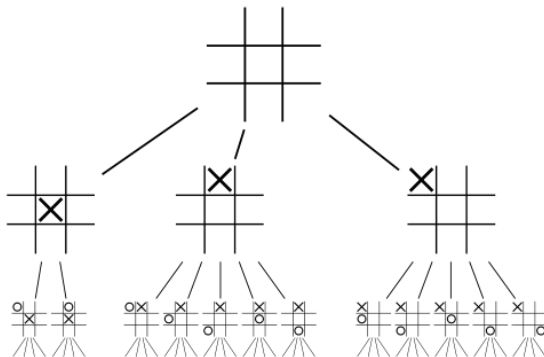
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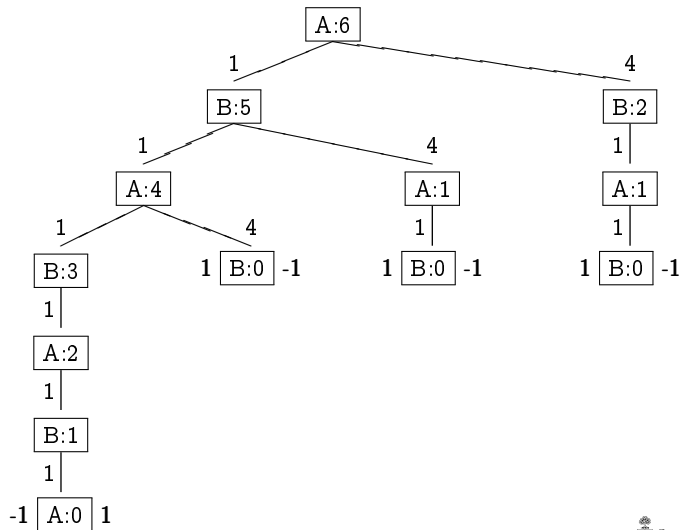

## minimax AKA negamax

This is a strong strategy that assumes both players are completely rational and have full information. They look at the consequences of each possible move and select the one with the best outcome for them. A very simplified diagram:



# minimax on subtract square

Subtract square doesn't spread out so fast:



## what about rough\_outcome?

Make an educated guess at the score without looking ahead any moves. In subtract square, this might be: “Can I win this move?” versus “Will I guarantee my opponent can win next move?” versus “Neither of the above”.

```
if is_pos_square(self.current_total):
    return SubtractSquareState.WIN
elif all([is_pos_square(self.current_total - n**2)
          for n in range(1, self.current_total + 1)
          if n**2 < self.current_total]):
    return SubtractSquareState.LOSE
else:
    return SubtractSquareState.DRAW
```



# BTNode

Change our generic **Tree** design so that we have two named children, **left** and **right**, and can represent an empty tree with **None**

```
class BTNode:
    '''Binary Tree node.'''

    def __init__(self, data, left=None, right=None):
        ''' (BTNode, object, BTNode, BTNode) -> NoneType

        Create BTNode (self) with data
        and children left and right.
        '''
        self.data, self.left, self.right = data, left, right
```





## special methods...

We'll want the standard special methods:

- ▶ `__eq__`

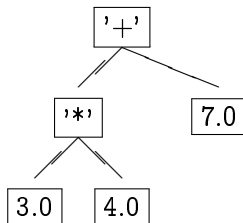
- ▶ `__str__`

- ▶ `__repr__`



# arithmetic expression trees

Binary arithmetic expressions can be represented as binary trees:





# inorder

A recursive definition:

- ▶ visit the left subtree **inorder**
- ▶ visit this node itself
- ▶ visit the right subtree **inorder**

The code is almost identical to the definition.



# preorder

- ▶ visit this node itself
- ▶ visit the left subtree in **preorder**
- ▶ visit the right subtree in **preorder**



# postorder

- ▶ visit the left subtree in **postorder**
- ▶ visit the rightsubtree in **postorder**
- ▶ visit this node itself



# definition

Add ordering conditions to a binary tree:

- ▶ data are comparable
- ▶ data in left subtree are less than node.data
- ▶ data in right subtree are more than node.data



# why binary search trees?

Searchs that are directed along a single path are efficient:

- ▶ a BST with 1 one has height 1
- ▶ a BST with 3 nodes may have height 2
- ▶ a BST with 7 nodes may have height 3
- ▶ a BST with 15 nodes may have height 4
- ▶ a BST with  $n$  nodes may have height  $\lceil \lg n \rceil$ .





## bst\_contains

If node is the root of a “balanced” BST, then we can check whether an element is present in about  $\lg n$  node accesses.

```
def bst_contains(node, value):  
    ''' (BTreeNode, object) -> value
```

Return whether tree rooted at node contains value.

Assume: node is the root of a BST.

```
>>> bst_contains(None, 5)  
False  
>>> bst_contains(BTreeNode(7, BTreeNode(5), BTreeNode(9)), 5)  
True  
, , ,  
  
# Use BST property to avoid  
# examining unnecessary nodes.
```