# CSC165 fall 2019 <br> end induction... <br> ...begin algorithm analysis 

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Using Course notes: more Induction

## Outline

notes

## Compure Sisere

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$\square$

## every set with $n$ elements has $2^{n}$ subsets

order of introductions... and intuition


## taking binary representation apart

suppose $n$ is a natural number with binary representation:

$$
n=\Sigma_{i=0}^{i=p} b_{i} 2^{i}=b_{p} 2^{p}+b_{p-1} 2^{p-1}+\cdots+2^{1} b_{1}+2^{0} b_{0}
$$

figure out $b_{0}$ based on whether $n$ is odd or even... suggesting

```
def natural_to_binary(n: int) -> str:
    # convert n to equivalent binary string
    bs = str (n % 2)
    n = n // 2
    while n > 0:
        bs = str (n % 2) + bs
        n = n // 2
```

    return bs
    
## time resource

How much time does this take?

```
def f(list_):
    for i in list_:
        print(i)
```


## assumptions, assumptions...

- "steps"
- ignore constant factors
- ignore "noise" for small input

We care about growth rate of time consumption

## formalizing assumptions

- $f$ absolutely dominates $g$
- $f$ dominates $g$ up to a constant factor
- $f$ eventually dominates $g$ up to a constant factor

What should domain and range of $f, g$ be?

## big-Oh, big-Omega, big-Theta

... and you're started on the Greek alphabet...

$$
f \in \mathcal{O}(g): \exists c, n_{0} \in \mathbb{R}^{+}, \forall n \in \mathbb{N}, n \geq n_{0} \Rightarrow f(n) \leq c g(n)
$$

try to show that $\forall a, b \in \mathbb{R}^{+}$, $a n+b \in \mathcal{O}\left(n^{2}\right)$

## $\forall a, b \in \mathbb{R}^{+}, a n+b \in \mathcal{O}\left(n^{2}\right)$

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```
big-Oh hierarchy when 0<a<b
recall \mathcal{O}(f)={g|g:\mathbb{N}->\mp@subsup{\mathbb{R}}{}{+}\wedge\existsc,\mp@subsup{n}{0}{}\in\mp@subsup{\mathbb{R}}{}{+},\foralln\in\mathbb{N},n\geq\mp@subsup{n}{0}{}=>g(n)\leqcf(n)}
\mp@subsup{log}{a}{}n versus log
n}\mp@subsup{}{}{a}\mathrm{ versus }\mp@subsup{n}{}{b}\mathrm{ (polynomial)
a}\mp@subsup{}{}{n}\mathrm{ versus }\mp@subsup{b}{}{n}\mathrm{ (exponential)
\mp@subsup{\operatorname{log}}{a}{}n versus n
n}\mp@subsup{}{}{a}\mathrm{ versus b }\mp@subsup{}{}{n
explore!
```


## properties

- reflexivity
- transitivity of big-Oh
- not symmetry (anti-symmetry...)


## products and sums

- $a f$
- $f \cdot g$
- $f+g$

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