CSC165 fall 2019 graph connectivity

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Using Course notes: graphs

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must be connected

 $\forall n \in \mathbb{N}, \exists M \in \mathbb{N}, \forall G = (V, E), (|V| = n \land |E| \ge M) \Rightarrow G \text{ is connected}?$



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maybe connected

 $\forall n \in \mathbb{N}, \exists G = (V, E), |V| = n \land |E| = n - 1 \land G \text{ is connected}$



must be disconnected

 $\forall n \in \mathbb{N}, \forall G = (V, E), (|V| = n \land |E| \le n - 2) \Rightarrow G \text{ is not connected}$

steps:

- natural to reason by removing an edge from a connected graph with n 1 edges...
- first need some results about which components of connected graphs have redundant edges (cycles)...
- then need some results about connected graphs without cycles (trees)...
- then reason about reducing an arbitrary connected graph to a tree...

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whew!

cycle

consecutively adjacent vertices $v_0, \ldots, v_k \in V \land k \ge 3$, all distinct except $v_0 = v_k$ $\forall G = (V, E), \forall e \in E, G \text{ connected } \Rightarrow (e \text{ in a cycle of } G \Leftrightarrow G - e \text{ connected })$



tree: connected, acyclic graph

removing any edge from a tree disconnects it $\forall G = (V, E), G$ is a tree $\Rightarrow |E| = |V| - 1$ but first...

 $\forall G = (V, E), (G \text{ is a tree } \land |V| \ge 2) \Rightarrow (\exists v \in V, d(v) = 1)$



main result...

 $\forall n \in \mathbb{N}^+, \forall G = (V, E), (G \text{ is a tree } \land |V| = n) \Rightarrow |E| = |V| - 1$



big picture...

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Notes

