

CSC236 fall 2012

subtle induction

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Using [Introduction to the Theory of Computation](#),
Section 1.2–1.3

Well-ordering example

$\forall n, m \in \mathbb{N}, n \neq 0, R = \{r \in \mathbb{N} \mid \exists q \in \mathbb{N}, m = qn + r\}$ has a smallest element

This is the main part of proving the existence of a unique quotient and remainder:

$$\forall m \in \mathbb{N}, \forall n \in \mathbb{N} - \{0\}, \exists q, r \in \mathbb{N}, m = qn + r \wedge 0 \leq r < n$$

The course notes use Mathematical Induction. Well-ordering is shorter and clearer.

Principle of well-ordering

Every non-empty subset of \mathbb{N} has a smallest element

Is there something similar for \mathbb{Q} or \mathbb{R} ?

For a given pair of natural numbers $m, n \neq 0$ does the set R satisfy the conditions for well-ordering?

$$R = \{r \in \mathbb{N} \mid \exists q \in \mathbb{N}, m = qn + r\}$$

If so, we still need to be sure that

1. $0 \leq r < n$
2. That q and r are unique — no other natural numbers would work

...in order to have

$$\forall m \in \mathbb{N}, \forall n \in \mathbb{N} - \{0\}, \exists q, r \in \mathbb{N}, m = qn + r \wedge 0 \leq r < n$$

Every non-empty subset of \mathbb{N} has a smallest element

$$\forall m \in \mathbb{N}, \forall n \in \mathbb{N} - \{0\}, \exists q, r \in \mathbb{N}, m = qn + r \wedge 0 \leq r < n$$



Every non-empty subset of \mathbb{N} has a smallest element

$$\forall m \in \mathbb{N}, \forall n \in \mathbb{N} - \{0\}, \exists q, r \in \mathbb{N}, m = qn + r \wedge 0 \leq r < n$$



Every non-empty subset of \mathbb{N} has a smallest element

$P(n)$: Every round-robin tournament with n players that has a cycle has a 3-cycle

Claim: $\forall n \in \mathbb{N} - \{0, 1, 2\}, P(n)$.

If there is a cycle $p_1 > p_2 > p_3 \dots > p_n > p_1$, can you find a shorter one?

Every non-empty subset of \mathbb{N} has a smallest element

$P(n)$: Every round-robin tournament with n players that has a cycle has a 3-cycle

Claim: $\forall n \in \mathbb{N} - \{0, 1, 2\}, P(n)$.

$$2^n > 10n : P(n)$$

Where do we start?

$$P(0) \checkmark$$

$$P(3) \times$$

$$P(7) \checkmark$$

$$P(1) \times$$

$$P(4) \times$$

$$P(8) \checkmark$$

$$P(2) \times$$

$$P(5) \times$$

$$\boxed{P(6) \checkmark} \text{ Base?}$$

It's not true for several low values of n . You could re-write the predicate as $P'(n) : 2^{n+6} > 10(n+6)$, but why not just start later?

omit Base case (for now) ← start at 6.

Assume $n \in \mathbb{N} - \{0, 1, 2, 3, 4, 5, ?\}$, and that $P(n)$

to trial, $2^n > 10n$ (IH)

$$\begin{aligned} \text{Then, } 2^{n+1} &= 2 \cdot 2^n > 20n \quad \# \text{ (by IH)} \\ &= 10n + 10n \\ &\geq \end{aligned}$$

$$10(n+1) = 10n + 10. \quad \# \text{ since } n \geq 6 \geq 1$$

So, $P(n+1)$ follows

So $\forall n \in \mathbb{N} - \{0, 1, 2, 3, 4, 5\}, P(n) \Rightarrow P(n+1)$

Conclude $\forall n \in \mathbb{N} - \{0, 1, 2, 3, 4, 5\}, P(n)$



$$2^n > 10n$$

Where do we start?

$$3^n \geq n^3 : P(n)$$

Check your induction step

$$P(0) \checkmark$$

$$P(1) \checkmark$$

$$P(2) \checkmark$$

$$P(3) \checkmark$$

$$P(4) \checkmark$$

$$P(5) \checkmark$$

$$P(6) \checkmark$$

Look at the graph.

$$3^n \geq n^3 : P(n)$$

Check your induction step

Omit base(s)

0, 1, 2, 3 Verify these separately,
ie they are base cases.

Assume $n \in \mathbb{N} - \{0, 1, 2\}$ and that $P(n)$ is true, that

$$\text{is } 3^n \geq n^3.$$

$$\text{Then } 3^{n+1} = 3 \cdot 3^n \geq 3n^3$$

$$3n^3 \geq (n+1)^3$$

$$= n^3 + n^3 + n^3$$

$$\geq \dots$$

$$n^3 + 3n^2 + 3n + 1$$

$$= (n+1)^3$$

since $n^3 = n^3$
and $n^3 \geq 3n^2$
and $n^3 \geq 3n + 1$
when $n \geq 3$

0 X
1 X
1 2 X
1 1
1 2 1
1 3 3 1

Notes