A2 - due N 10 days prove/disprove. CSC165 winter 2013 Mathematical expression 1×3, 3×5, 5×7, Danny Heap 1+2, 2+3, 3+4,4+5 heap@cs.toronto.edu 9×15 BA4270 (behind elevators) http://www.cdf.toronto.edu/~heap/165/W13/ 416-978-5899

Course notes, chapter 3

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## Outline

## notes



proof bout limits  

$$\forall e \in \mathbb{R}^+$$
  $|y \in \mathbb{R}| ||u - \pi| < \delta \Rightarrow |y^2 - \pi^2| < \varepsilon$   
Assume  $\varepsilon \in \mathbb{R}^+$  # to introduce  $\forall$   
 $\forall e \in \mathbb{R}^+$  # to introduce  $\forall$   
 $Assume |y - \pi| < \delta$  # to introduce  $\forall$   
 $Assume |y - \pi| < \delta \Rightarrow |y^2 - \pi^2| < \varepsilon$   
Then  $|y - \pi| < \delta \Rightarrow |y^2 - \pi^2| < \varepsilon$  # introduced  $\forall$   
 $Then \forall y \in \mathbb{R}, |y - \pi| < \delta \Rightarrow |y^2 - \pi^2| < \varepsilon$  # introduced  $\forall$   
 $Then \exists \delta \in \mathbb{R}^+$   $\forall y \in \mathbb{R}, |y - \pi| < \delta \Rightarrow |y^2 - \pi^2| < \varepsilon$   
# introduced  $\exists$   
 $Conclude \forall \varepsilon \in \mathbb{R}^+$   $\exists \delta \in \mathbb{R}^+, \forall y \in \mathbb{R}, |y - \pi| < \delta \Rightarrow |y^2 - \pi^2| < \varepsilon$   
# introduced  $\exists$ 

fill in the: Then  $|y^2 - \pi^2| = |y - \pi| \cdot |y + \pi| \#$  factor. < 5 | y+7 | # | y-7 | < 5, assump  $= \delta |y - \pi + \pi + \pi| \# \text{ and } 0$  $\leq \delta ||y - \pi| + 2\pi | \# \triangle - \text{inequally}$ # |-1+5|  $\angle S | S + 27 | \# \leq |1-1|+5 |$ 4 δ | 1 + 27 | # 8≤1 = S( 1+ 27) # 1, 27 +- pre  $\# S \leq \frac{\varepsilon}{1+2\pi}$  $# S = min(1) \frac{\varepsilon}{H21} works.$ 

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Sometimes your argument has to split to take into account possible properties of your generic element:

 $orall n \in \mathbb{N}, n^2+n$  is even

A natural approach is to factor  $n^2 + n$  as n(n + 1), and then consider the case where n is odd, then the case where n is even.



scratch assume 
$$n \in \mathbb{N}$$
 # in rider to introduce  $\forall$   
Then  $(\exists k \in \mathbb{N}, n = 2k) \vee (\exists k \in \mathbb{N}, n = 2k+1)$   
Case 1 assume  $\exists k \in \mathbb{N}, n = 2k$   
Then  $n^{2}+n = n(n+1)$  # factor.  
 $= 2k(2k+1) \#$  sub  $n = 2k$ .  
 $= 2(k(2k+1)) \#$  factoring.  
Then  $\exists k' \in \mathbb{N}, n^{2}+n = 2k'$   
 $\# k' = k(2k+1) \in \mathbb{N}, \text{ since } 2, 1, k \in \mathbb{N}$   
Then  $n^{2}+n$  is even.  
(ase 2 assume  $\exists k \in \mathbb{N}, n = 2k+1 \notin$   
Then  $n^{2}+n = n(n+1)$   
 $= (2k+1)(2k+1+1) \#$  sub  $\bar{m}$   
 $= (2k+1)(2k+1+1) \#$  sub  $\bar{m}$   
 $= (2k+1)(2k+2)$   
 $\exists imile$   
Since  $n^{2}+n$  is even in both possible ease,

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## get wrong right

Be careful proving a claim false. Consider the claim, for some suitably defined X, Y and P, Q:  $\chi = \gamma' = \xi / \xi$  $\forall x \in X, \forall y \in Y, P(x, y) \Rightarrow Q(x, y) \otimes (x, y) : \chi > y$  $\exists \chi \in \chi$ 

To disprove S, should you prove:

$$\bigvee \qquad orall x \in X, orall y \in Y, P(x,y) \Rightarrow 
eg Q(x,y)$$

What about  

$$\begin{array}{c}
 & P(x,y) \wedge^{\neg Q}(x,y) \\
 & X \in X, \forall y \in Y, \neg (P(x,y) \Rightarrow Q(x,y)) \\
 & X = Y = \mathbb{N} \quad P(x,y) \Rightarrow X \langle y \\
 & X \langle y \rangle & X \langle y \\
\end{array}$$
Explain why, or why not.

Define T(n) by:

$$\forall n \in \mathbb{N}$$
  $T(n) \Leftrightarrow \exists i \in \mathbb{N}, n = 7i + 1.$ 

Take some scrap paper, don't write your name on it, and fill in as much of the proof of the following claim as possible:

$$orall n \in \mathbb{N}, \, T(n) \Rightarrow \, T(n^2)$$

Now fill in as much of the **disproof** of the following claim as possible:

$$orall n \in \mathbb{N}, \, T(n^2) \Rightarrow T(n)$$

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Notes

