

CSC148 fall 2013

sorting big-oh
week 9

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Outline

more big-oh

running time analysis

algorithm's behaviour over large input (size n) is common way to compare performance

constant: $c \in \mathbb{R}^+$ (some positive number)

logarithmic: $c \log n$

linear: cn (probably not the same c)

quadratic: cn^2

cubic: cn^3

exponential: $c2^n$

horrible: cn^n or $cn!$

case: $\lg n$

this is the number of times you can divide n in half before reaching 1.

- ▶ refresher: $a^b = c$ means $\log_a c = b$.
- ▶ this runtime behaviour often occurs when we “divide and conquer” a problem (e.g. binary search)
- ▶ we usually assume $\lg n$ (log base 2), but the difference is only a constant:

$$2^{\log_2 n} = n = 10^{\log_{10} n} \implies \log_2 n = \log_2 10 \times \log_{10} n$$

- ▶ so we just say $\mathcal{O}(\lg n)$.

hierarchy

Since big-oh is an **upper-bound** the various classes fit into a hierarchy:

$$\mathcal{O}(1) \subseteq \mathcal{O}(\lg n) \subseteq \mathcal{O}(n) \subseteq \mathcal{O}(n^2) \subseteq \mathcal{O}(n^3) \subseteq \mathcal{O}(2^n) \subseteq \mathcal{O}(n^n)$$

scaling:

How well do these various sorts perform as the size of the problem (list length) increases? Time and compare.

term test #2

- ▶ Same time as lecture, but in EX300 (A–L) and EX310 (M–Z)
- ▶ You are responsible for lecture examples, assignment 2, labs, and exercises. That's where I'll look for suitable questions.
- ▶ Topics include, but not limited to: binary trees, linked lists, binary search trees, recursion on trees and lists, big-oh analysis