CSC236 fall 2012

structural induction

Danny Heap
heap@cs.toronto.edu
BA4270 (behind elevators)

http://www.cdf.toronto.edu/~heap/236/F12/
416-978-5899

Using Introduction to the Theory of Computation,
Chapter 4, Section 1.1
Outline

Equivalence of inductions

Structural induction

Notes
The cycle is proved in the text, here is one link. Suppose you believe MI, and you have shown for some property $P$:

$$\forall n \in \mathbb{N}, (\forall 0 \leq i < n, P(i)) \Rightarrow P(n) \quad (1)$$

Now define a slightly different predicate:

$P'(n) : \forall 0 \leq i \leq n, P(i)$, in other words, $P(i)$ is true up to and including $n$. Using only MI prove $\forall n, P'(n)$:

**Base case:** Since we showed (1), and there are no natural numbers smaller than 0, we have $P'(0)$.

**Induction step:** Assume $n$ is an arbitrary natural number and that $P'(n)$ is true. It follows from (1) that $P(n + 1)$ is true, and hence $P'(n + 1)$ is true.
Define sets inductively

...so as to use induction on them later

One way to define the natural numbers:

$\mathbb{N}$: The smallest set such that

1. $0 \in \mathbb{N}$
2. $n \in \mathbb{N} \Rightarrow n + 1 \in \mathbb{N}$.

By smallest I mean $\mathbb{N}$ has no proper subsets that satisfy these conditions. If I leave out smallest, what other sets satisfy the definition?

$\downarrow \mathbb{Q}, \mathbb{R}, \mathbb{C}$,
What can you do with it?

The definition on the previous page defined the simplest natural number (0) and the rule to produce new natural numbers from old (add 1). Proof using Mathematical Induction work by showing that 0 has some property, and then that the rule to produce natural numbers preserves the property, that is

1. Prove $P(0)$
2. Prove that $\forall n \in \mathbb{N}, \ P(n) \Rightarrow P(n+1)$. 
Other structurally-defined sets

\[ E \text{ includes } x, y, z, (x+y), (y+z), (y+x), (y \times x), ((x+y) \times (y+z)) \]

Define \( E \): The smallest set such that

- \( x, y, z \in E \)
- \( e_1, e_2 \in E \Rightarrow (e_1 + e_2), (e_1 - e_2), (e_1 \times e_2), \) and \((e_1 \div e_2) \in E.\)

Form some expressions in \( E \). Count the number of variables (symbols from \( \{x, y, z\}\)) and the number of operators (symbols from \( \{+, \times, \div, -\}\)). Make a conjecture.

\[ \forall e \in E \quad \text{var}(e) = \text{op}(e) + 1 \]
Structural induction

\[ P(e) : \text{vr}(e) = \text{op}(e) + 1 \]

\[(x + y) + z\]

To prove that a property is true for all \( e \in \mathcal{E} \), parallel the recursive set definition:

- **Base case:** Show that the property is true for the simplest members, \( \{x, y, z\} \)
- **Induction step:** Show “inheritance”: if \( P(e_1) \) and \( P(e_2) \), then all possible combinations \( (e_1 + e_2) \), \( (e_1 - e_2) \), \( (e_1 \times e_2) \), and \( (e_1 \div e_2) \) have the property.

Conclude that the property is true of all elements of \( \mathcal{E} \).
**Structural induction**

\[ P(e) : \text{vr}(e) = \text{op}(e) + 1 \]

# Variables \# operators

Prove \( \forall e \in \mathcal{E}, P(e) \)

**Proof (structural induction).**

**Base case** for basic elements \( x, y, z \in \mathcal{E}, \) each has \( \text{vr}(e) = 1 = 0 + 1 = \text{op}(e) + 1, \) so claim holds base case.

**Induction step** Assume \( P(e_1) \) and \( P(e_2) \) hold for some arbitrary \( e_1, e_2 \in \mathcal{E}. \) Denote \( \circ \in \mathcal{E}_\circ , \) \( x_1 \in \mathcal{X}, \) and note that \( \text{vr}((e_1 \circ e_2)) = \text{vr}(e_1) + \text{vr}(e_2). \) Also note \( \text{op}((e_1 \circ e_2)) = \text{op}(e_1) + \text{op}(e_2) + 1. \) That means \( \text{op}((e_1 \circ e_2)) = \text{op}(e_1) + \text{op}(e_2) + 1 \)

\[ \begin{array}{c}
\text{IH} \\
= \text{vr}(e_1) - 1 + \text{vr}(e_2) - 1 + 1 \\
= \text{vr}(e_1) + \text{vr}(e_2) - 1 = \text{vr}((e_1 \circ e_2)) - 1,
\end{array} \]

**We have shown** \( P((e_1 \circ e_2)) \)
Recursive definition

Fibonacci sequence

This sequence comes up in applied rabbit breeding, the height of AVL trees, and the complexity of Euclid’s algorithm for the GCD:

\[
F(n) = \begin{cases} 
n & n < 2 \\
F(n - 2) + F(n - 1) & n \geq 2 
\end{cases}
\]

What is the sum of \(n\) Fibonacci numbers?
Fibonacci numbers

What is \( \sum_{i=0}^{n} F(i) \)?

Claim: \( \forall n \in \mathbb{N}, P(n) \).

Base case

Induction step: Assume \( n \) is an arbitrary natural number and that \( P(n) \) is true.

Then,

\[
\sum_{i=0}^{n+1} F(i) = \left( \sum_{i=0}^{n} F(i) \right) + F(n+1)
\]

\[
= F(n+2) - 1 + F(n+1) \quad \text{(by IH)}
\]

\[
= F(n+3) - 1, \quad \text{which is } P(n+1) \quad \text{since } n+3 \geq 2
\]

Therefore, \( \forall n \in \mathbb{N}, P(n) \Rightarrow P(n+1) \).

Conclude: \( \forall n \in \mathbb{N}, P(n) \).