CSC 165
proof by cases
week 7, lecture 3
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Assignment 1 graded, pick up at end of lecture
T2 I have some.
Assignment 2 posted, due November 13th.
Tutorials - all about proof structure
proof by cases

You can prove by induction (CSC236) that:

\[ \forall m \in \mathbb{N}, \forall n \in \mathbb{N}, n \neq 0 \Rightarrow \exists q \in \mathbb{N}, \exists r \in \mathbb{N}, m = qn + r \text{ and } n > r \geq 0 \]

\( \exists! \) is a compact way of saying there exists exactly one.

- \( q \) and \( r \) are called the quotient and remainder, respectively. We also denote \( r \) by \( m \mod n = r \)

A consequence is that any natural number \( n \) has a remainder of either 0, 1, or 2 after division by 3. What possible remainders are there for perfect squares after division by 3?

- \( 0^2 \mod 3 = 0 \)
- \( 1^2 \mod 3 = 1 \)
- \( 2^2 \mod 3 = 1 \)
- \( 3^2 \mod 3 = 0 \)
- \( 4^2 \mod 3 = 1 \)
- \( 5^2 \mod 3 = 1 \)
- \( 6^2 \mod 3 = 0 \)
Prove: $\forall n \in \mathbb{N}, n^2 \mod 3 \neq 2$

I think you’ll need cases for different possible results of $n \mod 3$

Assume $n \in \mathbb{N}$ is generic.

Then $n \mod 3 \in \{0, 1, 2\}$ # defn of $a \mod b$

Case 1: assume $n \mod 3 = 0$

Then $\exists q' \in \mathbb{N}, n = 3q' + 0$ # defn of $n \mod 3 = 0$

Pick $q' \in \mathbb{N}, n = 3q'$ # since it exists.

Then $n^2 = 3(3q'^2)$ # algebra.

Then $\exists q' \in \mathbb{N}, n^2 = 3q' + 0$ # $3q'^2$ since $3, q' \in \mathbb{N}$ and # $\mathbb{N}$ is closed under $*$

Then $n^2 \mod 3 = 0 \neq 2$

Case 2: assume $n \mod 3 = 1$

Then $\exists q' \in \mathbb{N}, n = 3q' + 1$ # defn of $n \mod 3 = 1$

Then $n^2 \mod 3 = 1 \neq 2$
Case 3, assume \( n \mod 3 = 2 \)

Then \( \exists q' \in \mathbb{N}, n = 3q' + 2 \)

Pick \( q \in \mathbb{N}, n = 3q + 2 \)

Then \( n^2 = 9q^2 + 12q + 4 \)

\[ = 3(3q^2 + 4q + 1) + 1 \]

So \( \exists q' \in \mathbb{N}, n^2 = 3q' + 1 \)

Then \( n^2 \mod 3 = 1 \neq 2 \)

Conclude \( \forall n \in \mathbb{N}, n^2 \mod 3 \neq 2 \) since true in all possible cases.
Performing arithmetic on integers, and then taking the remainder, can be interchanged with taking the remainder and then performing the arithmetic. Number theory and cryptography use many such techniques.

Example: $\forall m \in \mathbb{N}, \forall n \in \mathbb{N}, (m \mod 7 = 3 \land n \mod 7 = 4) \Rightarrow mn \mod 7 = 5$

An important step is to remember that $m \mod 7 = 3$ means there are natural numbers $q$ and $r$ such that $m = 7q + 3$. 
Prove: \( \forall m \in \mathbb{N}, \forall n \in \mathbb{N}, (m \mod 7 = 3 \land n \mod 7 = 4) \Rightarrow mn \mod 7 = 5. \)

Assume \( m, n \in \mathbb{N} \) \# generic pair in \( \mathbb{N} \).

Assume \( (m \mod 7 = 3 \land n \mod 7 = 4) \) \# antecedent.
Then $mn \mod 7 = 5$

Conclude $\forall m, n \in \mathbb{N}, (m \mod 7 = 3 \land n \mod 7 = 4) \Rightarrow mn \mod 7 = 5$
structural transformations

Pages 48–49 of the notes list some standard transformations used in proofs. These have the virtue that many proof practitioners will recognize them when you use them.

Work back over proofs to find:

- implication introduction
- universal introduction
- existential introduction
- existential elimination