Question 1. [4 marks]

Consider the following schema for a database about transit routes, stops along the routes, and the scheduled time to reach the stops on a route (for various trips along the route).

Relations

- Route(name, num, mode)
- Stop(route, stopNum, location)
- StopTime(route, stopNum, tripNum, time)

Integrity constraints

- $\Pi_{mode} Route \subseteq \{ "subway", "streetcar", "bus", "LRT" \}$
- Stop[route] $\subseteq$ Route[num]
- StopTime[route, stopNum] $\subseteq$ Stop[route, stopNum]

Which of the following statements are enforced by the schema? Circle one answer for each. If the statement is enforced, say what part of the schema enforces it. If it is not enforced, write an integrity constraint that would enforce it, using one of the two forms defined in the textbook.
(No penalty for wrong answers, but no marks without a correct explanation or integrity constraint.)

1. Two routes cannot have the same name.

   Enforced This part of the schema enforces it:

   ![Not enforced](image)

   This new integrity constraint would enforce it:

   $\sigma_{one.num \neq two.num \land one.name = two.name}(\rho_{one} Route \times \rho_{two} Route) = \emptyset$

2. Every value for route in relation StopTime occurs as a value for num in relation Route.

   Enforced This part of the schema enforces it:

   ![Not enforced](image)

   This new integrity constraint would enforce it:

   Every route in StopTime must occur with a stopNum.
   Because StopTime[route, stopNum] $\subseteq$ Stop[route, stopNum],
   each of these pairs must occur in Stop, so certainly the route occurs in Stop.
   And because Stop[route] $\subseteq$ Route[num], the route appears as num in relation Route.
Question 2. [4 marks]

Suppose we have two relations: Patient(PID, height) and Caresfor(PID, doctor). Consider the following instance of that schema:

<table>
<thead>
<tr>
<th>Patient</th>
<th>PID</th>
<th>height</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>12</td>
<td>11</td>
</tr>
<tr>
<td></td>
<td>25</td>
<td>12</td>
</tr>
<tr>
<td></td>
<td>96</td>
<td>11</td>
</tr>
<tr>
<td></td>
<td>20</td>
<td>12</td>
</tr>
<tr>
<td></td>
<td>33</td>
<td>33</td>
</tr>
<tr>
<td></td>
<td>44</td>
<td>8</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Caresfor</th>
<th>PID</th>
<th>doctor</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>33</td>
<td>44</td>
</tr>
<tr>
<td></td>
<td>20</td>
<td>44</td>
</tr>
<tr>
<td></td>
<td>25</td>
<td>30</td>
</tr>
<tr>
<td></td>
<td>12</td>
<td>30</td>
</tr>
<tr>
<td></td>
<td>20</td>
<td>18</td>
</tr>
<tr>
<td></td>
<td>96</td>
<td>30</td>
</tr>
<tr>
<td></td>
<td>12</td>
<td>20</td>
</tr>
</tbody>
</table>

Part (a) [2 marks]

Give the result (schema and data) returned by the following query. Use the same tabular format as above.

Temp(PID, height) := Π_{P1.PID,P1.height}(σ_{P1.height>P2.height}(ρ_{P1}(Patient) × ρ_{P2}(Patient)))

Temp2(PID, height) := Π_{T1.PID,T1.height}(σ_{T1.height>T2.height}(ρ_{T1}(Temp) × ρ_{T2}(Temp)))

Π_{PID,doctor}(Temp − Temp2) ⊗◁ Caresfor

Solution:

<table>
<thead>
<tr>
<th>PID</th>
<th>doctor</th>
</tr>
</thead>
<tbody>
<tr>
<td>12</td>
<td>30</td>
</tr>
<tr>
<td>12</td>
<td>20</td>
</tr>
<tr>
<td>96</td>
<td>30</td>
</tr>
</tbody>
</table>

Part (b) [2 marks]

Describe what this query computes. Do not describe the steps it takes, only what is in the result, and make your answer general to any instance of the schema.

Solution:

It finds the second-shortest patient(s) and for each one of them, reports the PID and all of their doctors.
Question 3. [15 marks]

We used the following schema many times in lecture:

**Relations**
- Students(SID, surName, campus)
- Courses(CID, cName, WR)
- Offerings(OID, CID, term, instructor)
- Took(SID, OID, grade)

**Integrity constraints**
- Offerings[CID] ⊆ Courses[CID]
- Took[SID] ⊆ Students[SID]
- Took[OID] ⊆ Offerings[OID]

**Part (a) [6 marks]**

Write a query in relational algebra to report the surname and SID of all students who have taken at most one course in which no one in that offering of the course earned a grade of 100. Use only the basic operators \(\Pi, \sigma, \bowtie, \times, \cap, \cup, -, \rho\).

**Solution:**

\[
\begin{align*}
Bad(OID) &:= (\Pi_{OID}Offering) - (\Pi_{OID}\sigma_{grade=100}Took) \\
TookBad(SID, OID) &:= \Pi_{SID,OID}(Took \bowtie \bowtie Bad) \\
TwoPlus(SID) &:= \sigma_{TB_1.SID=TB_2.SID \land TB_1.OID \neq TB_2.OID}(\rho_{TB_1.TookBad \times \rho_{TB_2.TookBad}}) \\
\Pi_{SID}(Students) - TwoPlus
\end{align*}
\]

**Part (b) [1 mark]**

Can the following query be expressed using those same basic operators (\(\Pi, \sigma, \bowtie, \times, \cap, \cup, -, \rho\)):

Report the closest grade to 50 among the passing grades in any offering of csc363. Circle one answer.

(One mark for a correct answer and -0.5 for an incorrect answer.)

Yes No
Part (c) [8 marks]
Consider all the students who’ve ever taken csc333. Suppose we want to find the instructors who have given every one of them a passing grade in at least one course. (They need not have been in the same offering or course.) Which of the following syntactically legal queries will report that?
(2 marks for each correct answer, -1 for each incorrect answer.)

1. \( Takers(SID) := \Pi_{SID}\sigma_{CID = \text{“csc333”}} (\text{Took} \bowtie \text{Offering}) \)
   \( Did(SID, instructor) := \Pi_{SID, instructor}\sigma_{grade \geq 50}(\text{Took} \bowtie \sigma_{CID = \text{“csc333”}} \text{Offering}) \)
   \((\Pi_{instructor, Offering}) - (\Pi_{instructor}[Takers \bowtie \Pi_{instructor, Offering} - Did]) \)
   Correct \[\text{Incorrect}\]

2. \( Takers(SID) := \Pi_{SID}\sigma_{CID = \text{“csc333”}} (\text{Took} \bowtie \text{Offering}) \)
   \( Did(SID, instructor) := \Pi_{SID, instructor}\sigma_{grade \geq 50}(\text{Took} \bowtie \text{Offering}) \)
   \((\Pi_{instructor, Offering}) - (\Pi_{instructor}[Takers \bowtie \Pi_{instructor, Offering} - Did]) \)
   Correct \[\text{Incorrect}\]

3. \( Takers(SID) := \Pi_{SID}\sigma_{CID = \text{“csc333”}} (\text{Took} \bowtie \text{Offering}) \)
   \( Did(SID, instructor) := \Pi_{SID, instructor}(\sigma_{grade \geq 50}\text{Took} \bowtie [\Pi_{OID, instructor, Offering}]) \)
   \((\Pi_{instructor, Offering}) - (\Pi_{instructor}[Takers \bowtie \Pi_{instructor, Offering} - Did]) \)
   Correct \[\text{Incorrect}\]

4. \( Takers(SID) := \Pi_{SID}\sigma_{CID = \text{“csc333”}} (\text{Took} \bowtie \text{Offering}) \)
   \( Did(SID, instructor) := \Pi_{SID, instructor}(\sigma_{grade \geq 50}\text{Took} \bowtie [\Pi_{OID, instructor, Offering}]) \)
   \((\Pi_{instructor, Offering}) - (\Pi_{instructor}[Takers \bowtie \Pi_{instructor, Offering} - \Pi_{instructor, Did}]) \)
   Correct \[\text{Incorrect}\]
Question 4. [7 marks]

Consider the following schema about athletes and their results in the long jump event:

```
create table athlete (
    aID int primary key,
    name text not null
); create table longjump (
    who int,
    distance float not null,
    t timestamp,
    primary key (who, t),
    foreign key (who) references athlete(aID)
);```

Part (a) [3 marks]

Complete the `where` condition in the following SQL query so that it reports the aID and name of the athlete who jumped the greatest distance. If there is a tie, report them all. If there are duplicates, report them all.

```
select aid, name
from longjump natural join athlete
where
```

Solution:

Here are two correct answers:

```
aid = who and distance >= all (select distance from longjump);
```

and

```
aid = who and distance = (select max(distance) from longjump);
```

Part (b) [1 mark]

Could any query that begins as stated above actually produce duplicate tuples?
(One mark for a correct answer and -0.5 for an incorrect answer.)

[YES] No

Part (c) [3 marks]

Suppose we have a relation `R(a, b, c)`. Consider this query:

```
select _____________ from R group by b;
```

Which of the following could go in the select list for this query? Circle Okay or Error for each.
(One mark for each correct answer and -0.5 for each incorrect answer.)

<table>
<thead>
<tr>
<th>max(a)</th>
<th>Okay</th>
<th>Error</th>
</tr>
</thead>
<tbody>
<tr>
<td>c</td>
<td>Okay</td>
<td>Error</td>
</tr>
<tr>
<td>min(c)</td>
<td>Okay</td>
<td>Error</td>
</tr>
</tbody>
</table>