CSC148 winter 2018

linked lists, iteration, mutation  week 4

Danny Heap
heap@cs.toronto.edu / BA4270 (behind elevators)
http://www.teach.cs.toronto.edu/~csc148h/winter/
416-978-5899

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Outline

linked lists

mutation

linked list queues
why linked lists?

Regular Python lists are flexible and useful, but overkill in some situations. They allocate large blocks of contiguous memory, which becomes increasingly difficult as memory is in use.

Linked list nodes reserve just enough memory for the object value they want to refer to, a reference to it, and a reference to the next node in the list.
linked lists, two concepts

There are two useful, but different, ways of thinking of linked list nodes

1. as lists made up of an item (value) and a sub-list (rest)

2. as objects (nodes) with a value and a reference to other similar objects

For now, will take the second point-of-view, and design a separate “wrapper” to represent a linked list as a whole.
a node class

class LinkedListNode:
    
    Node to be used in linked list

    === Attributes ===
    next_ - successor to this LinkedListNode
    value: data this LinkedListNode represents
    
    def __init__(self, value: object, next_: Union["LinkedListNode", None]=None) -> None:
        
        Create LinkedListNode self with data value and successor next_.
        
        self.value, self.next_ = value, next_
a wrapper class for list

The list class keeps track of information about the entire list — such as its front, back, and size.

class LinkedList:
    
    Collection of LinkedListNodes

    === Attributes ==
    front - first node of this LinkedList
    back - last node of this LinkedList
    size - number of nodes in this LinkedList, >= 0
    
    front: Union[LinkedListNode, None]
    back: Union[LinkedListNode, None]
    size: int

    def __init__(self):
        
        Create an empty linked list.
        
        self.front, self.back, self.size = None, None, 0
division of labour

Three reasonable ways to define `__eq__`

Some of the work of special methods is done by the nodes:

- `__str__`
  
  \[
  \begin{array}{c}
  \text{ID = = ID} \\
  \text{Value = = Value} \\
  \text{Next = = next} \\
  \end{array}
  \]

- `__eq__`

Once these are done for nodes, it’s easy to do them for the entire list.
walking a list

Make a reference to (at least one) node, and move it along the list:

```python
cur_node = self.front
while <some condition here...>:
    # do something here...
    cur_node = cur_node.nxt
```

this name "moves" along the list... check for None!
prepend
create a new node and add it before self.front...
Check (possibly) every node

```python
cur_node = self.front
while <some condition here...>:
    # do something here...
    cur_node = cur_node.nxt
return False
```
Should enable things like

```python
>>> print(lnk[0])
5
```

... or even

```python
>>> print(lnk[0:3])
5 -> 4 -> 3 ->|
```

```
lnk.Size = 3
lnk[3] → IndexError
lnk[2]  
lnk[1] ✓
lnk[0] ✓
lnk[-1] ✓
lnk[-2]
lnk[-3]
lnk[-4]
```
append

We’ll need to change...

- last node
- former last node
- back
- size
- possibly front

draw pictures!!
We need to find the second last node. Walk two references along the list.

```python
prev_node, cur_node = None, lnk.front
# walk along until cur_node is lnk.back
while <some condition>:
    prev_node = cur_node
    cur_node = cur_node.nxt
```
something linked lists do better than lists?

list-based Queue has a problem: adding or removing will be slow.
symmetry with linked list

see previous!

which end of a linked list would be best to add, which to remove? why?? think about front, back...
build pop_front
  _ delete_front
  _ then pop_front

... already have append
revisit Queue API

these are now easy

use an underlying LinkedList
revisit Stack API while we’re at it

also use an underlying LinkedList
they’re all Containers

see our code

stress drive them through container_cycle, in container_timer.py:

- list-based Queue
- linked-list-based Queue
- list-based Stack
- linked-list-based Stack
what matters is the growth rate

when the list is 10x as big
what happens to run-time?

as Queue grows in size, list-based-Queue bogs down, becomes impossibly slow
notes...